

# Rapid Applications Prototyping in Haskell

## Workshop Session 4

2018-09-02

# Objectives

- ▶ Recap basic types and type classes.
- ▶ Using RIO and stack to write standalone scripts.
- ▶ More higher-kinded typeclasses, Foldable, Traversable, Applicative.
- ▶ Quick look at algebraic-graphs.

# Quick Recap

- ▶ **Values** are things we want to compute.
- ▶ **Types** are ways of classifying a range of values.
- ▶ **Functions** are ways of turning values of one type into values of another type
- ▶ **Type Classes** are ways of classfying types by an interface of functions available for that type.
- ▶ **Higher Kinded Types** make new types out of existing types or “functions at the type level”

## Simple scripts

Make a new blank file “foo.hs” and put

```
#!/usr/bin/env stack
-- stack runhaskell --resolver lts-12.13 --package rio

{-# LANGUAGE NoImplicitPrelude #-}
{-# LANGUAGE OverloadedStrings #-}

import RIO

main = runSimpleApp $ do
    logInfo $ "Hello World"
```

Now you can `chmod a+x foo.hs` and run it with `./foo.hs`.

Remember you can `export RIO_VERBOSE=true` to add color and timestamps.

## Shelling out

You can use the `proc` functiona to shell out.

```
import RIO
import RIO.Process

main = runSimpleApp $ do
  logInfo $ "Hello World"
  proc "touch" ["tiny kitten"] (runProcess_)
  (c, s, e) <- proc "uname" [] (readProcess)
  logInfo $ displayShow $ s
```

# Examples

This works really well with some libraries like diagrams, inline-r, clay and lucid.

Here are some links:

- ▶ [diagrams](#)
- ▶ [inline-r](#)

# Foldables

Foldables are things that can be iterated over and collapsed to a result.

Type :info Foldable

```
class Foldable (t :: * -> *) where
  foldr :: (a -> b -> b) -> b -> t a -> b
  ...
```

The first argument is a binary operation, like (+), (++) or (:)

The second argument is an initial value to start with (in case the Foldable is empty).

The third argument is the foldable, with some `as` inside.

# Foldable

Maybe and [] are both Foldable, so we can do things like:

```
sumList :: [Int] -> Int
sumList = foldr (+) 0
```

But why do this for just lists? We can sum any foldable with the same definition

```
sumFoldable = foldr (+) 0
```

What's the type of sumFoldable?

# DeriveFoldable

We can derive Foldable in the same way we derived Functor for well behaved data structures

```
{-# LANGUAGE DeriveFoldable #-}
```

```
{-# LANGUAGE DeriveFunctor  #-}
```

```
data V3 a = V3 a a a  
  deriving (Eq, Show, Functor, Foldable)
```

# Algebraic Graphs

algebraic-graphs is a great library that defines the following data structure.

```
data Graph a =  
    Empty  
  | Vertex a  
  | Connect (Graph a) (Graph a)  
  | Overlay (Graph a) (Graph a)  
deriving (Functor, Foldable, Show, Traversable)
```

Graph is higher kinded, takes a type parameter 'a', which allows us to talk about a “graph of 'a's”

This probably looks like a peculiar way to describe a graph. What values can a graph like this be?

Graphs are naturally Functors and Foldables.

# Algebraic Graphs

```
#!/usr/bin/env stack
{- stack runhaskell --resolver lts-12.13
   --package rio --package algebraic-graphs -}

{-# LANGUAGE NoImplicitPrelude #-}

import RIO
import Algebra.Graph

myGraph :: Graph String
myGraph = overlay
    (vertex "spatula")
    (connects $ [vertex "quux",
                 vertices ["foo", "bar", "baz"]])

main = runSimpleApp $ do
    logInfo $ displayShow $ myGraph
    logInfo $ displayShow $ foldr (++) "" myGraph
```

# Currying/Uncurrying

We know that haskell has two ways of representing two argument functions.

```
(a, b) -> c  
a -> b -> c
```

Haskell provides two functions to switch between these representations:

```
curry :: ((a, b) -> c) -> a -> b -> c  
uncurry :: (a -> b -> c) -> (a, b) -> c
```

This of course has the property that

```
curry . uncurry = id  
uncurry . curry = id
```

# Applicative Functors

A lax monoidal functor, or “applicative” functor is a functor that can lift functions of multiple arguments.

Remember the signature for `fmap` `fmap :: (a -> b) -> f a -> f b`

```
class Functor f => Applicative f where
  liftA2 :: (a -> b -> c) -> f a -> f b -> f c
  (<*>) :: f (a -> b) -> f a -> f b
```

`liftA2` is a two argument analogue of `fmap`, it lifts a two argument function into the higher kinded type.

The `(<*>)` is pronounced **splat**, sometimes called ‘ap’, it is defined as `(<*>) = liftA2 id`

`Maybe`, `IO`, `[]`, and the hom functor `(r ->)` are all `Applicative`s

# Applicative Functors

One project euler problem says

“Find the largest palindrome made from the product of two 3-digit numbers.”

What are all the products of two digit numbers?

We can use `liftA2` to take normal multiplication of numbers and turn it into a function that multiplies two lists of numbers.

```
liftA2 (*) [100..999] [100..999]
```

Sometimes the “fmap-then-splat” notation is used to the same effect.

```
(*) <$> [100..999] <*> [100..999]
```

Similarly we can lift binary operations to operations which act on two Maybes, or two IO values.

# Applicative Functors

Another list example:

```
data Value = Two | Three | Four | Five | Six |  
            Seven | Eight | Nine | Ten |  
            Jack | Queen | King | Ace  
    deriving (Eq, Ord, Enum, Show, Bounded)
```

```
data Suit = Hearts | Clubs | Spades | Diamonds  
    deriving (Eq, Ord, Enum, Show, Bounded)
```

```
type Card = (Value, Suit)
```

```
type Desk = [Card]
```

```
deck = liftA2 (,) [Two..] [Hearts..]  
-- deck = (,) <$> [Two..] <*> [Hearts..]
```

# The Hom Functor is Applicative

Let's continue the euler problem.

Recall the hom functor  $(r \rightarrow)$ , or “all of the functions emanating from a type  $r$ ”.

Just fill it in algebraically to see what `fmap` and `splat` do.

```
fmap :: (a -> b) -> f a -> f b
```

becomes

```
fmap :: (a -> b) -> (r -> a) -> (r -> b)
```

This is just the same as the dot `(.)` operator, so the `fmap` for the hom functor is just the dot operator `(.)`.

# The Hom Functor is Applicative

```
liftA2 :: a -> b -> c -> f a -> f b -> f c  
(<*>) :: f (a -> b) -> f a -> f b
```

becomes

```
liftA2 :: a -> b -> c -> (r -> a) -> (r -> b) -> (r -> c)  
(<*>) :: (r -> a -> b) -> (r -> a) -> r -> b
```

So what good is this?

# IsPalindrome

Let's try to make an `isPalindrome :: String -> Bool` that decides whether or not a string is a palindrome.

```
isPalindrome x = ???
```

# IsPalindrome

```
isPalindrome x = x == reverse x
```

Put this into <http://pointfree.io/> and see what happens.

# IsPalindrome

So equivalently:

```
isPalindrome = (==) <*> reverse
```

Can you say why?

# Traversable

A Traversable is just something that is both Functor and Foldable.

```
class (Functor t, Foldable t) => Traversable t
```

This has some interesting effects when you mix it with an applicative.

```
traverse    :: (Traversable t, Applicative f) =>  
              (a -> f b) -> t a -> f (t b)
```

```
sequenceA   :: (Traversable t, Applicative f) =>  
              t (f a) -> f (t a)
```

Functions of the form  $(a \rightarrow f\ b)$  for some function  $f$  are called “Stars”.

`traverse` takes a  $a \rightarrow f\ b$  and turns it into a  $t\ a \rightarrow f\ (t\ b)$

# Traversable

`sequenceA` is just equivalent to `traverse id`

There's also a function `for` which is just `traverse` with its arguments flipped, and works like a for loop.

```
for :: (Traversable t, Applicative f) => t a -> (a -> f b
```

We'll look more at these when we do monads, for now let's try a simple sequence.

# Traversable

Recall our graph earlier `myGraph :: Graph String`

Remember that a `String = [Char]`, so, `myGraph :: Graph [Char]`

`Graph` is a `Traversable`, and `[]` is an `Applicative`, so this is a `t (f a)`

So we can sequenceA `myGraph`

What happens?

# Traversable

It turns a `Graph [Char]` into a `[Graph Char]`, or a list of possible Graphs of different characters.

```
sequenceA myGraph  
:t sequenceA myGraph
```

How many Graphs has it produced?

```
length $ sequenceA myGraph
```